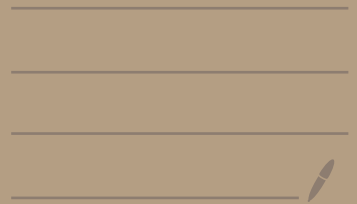


Exp. Deg. Lower bounds on

Invariant rings.

(Adv. Math, 2020 j/w Harm Derksen)



Motivation:

$$G \curvearrowright V$$

(gp reductive)
V rational

$$G = SL_n \times SL_n \times SL_n \curvearrowright V = \mathbb{C}^n \otimes \mathbb{C}^n \otimes \mathbb{C}^n$$

$\mathbb{C}[V]$ = polynomials in V

$$[V = \mathbb{C}^n, \mathbb{C}[V] = \mathbb{C}[x_1, \dots, x_n]]$$

$$\mathbb{C}[V]^G = \left\{ f \in \mathbb{C}[V] \mid \begin{array}{l} f(gv) = f(v) \\ \forall g \in G, v \in V \end{array} \right\}$$

Landmark Results:

① Invariant rings $(\mathbb{C}[V]^G)$ are
finitely generated.

Hilbert

Questions:

- ① What are the generators?
 - ①' Minimal set of generators
 - ② Relations among generators.
- } → 19th century

Realize over several decades →

- Can do this for some small cases
moderate cases

- It's pretty difficult in general.

→ ⇒ Complexity theory & Computers

Natural: How hard to compute?

[GCTV: Mulmuley.]

A naive measure of computational hardness is degree bounds.

Problem (Deg. bds)

Give strong bounds for

$$\beta(G, V) = \min \left\{ d \mid \mathbb{C}[V]^{\leq d} \text{ generate } \mathbb{C}[V]^G \right\}$$

Complexity aside \rightarrow deg d invariants are computable.

but without a degree bound, cannot get an algorithm naively to compute generators.

Hilbert: $\beta(G, V) < \infty$

Popov
80's: $\beta(G, V) \sim 2^{2^n}$

Derksen
2000: $\beta(G, V) \sim 2^{\text{poly}(n)}$

[Derksen, M. 2015]: L-R action

$SL_n \times SL_n$
" G
 $\hookrightarrow \text{Mat}_{n,n}^m$
" V

$[(\mathbb{C}^n \otimes \mathbb{C}^n)^{\oplus m}]$

$\beta(G, V) \leq n^6$

IQS \rightarrow

RIT $\in P$

det $\left(\begin{array}{|c|c|c|} \hline & & \\ \hline & & \\ \hline & & \\ \hline \end{array} \right)$

each block is a
lin. comb of $x_1 \rightarrow x_m$

L-R action \longrightarrow RIT

$G \curvearrowright V$ \rightsquigarrow Null Cone Membership (NCM)

Naive Hope: Better deg. bds \Rightarrow faster algos for NCM

Next generalization of the L-R action is

$$G = SL_n \times SL_n \times SL_n \curvearrowright V = (\mathbb{C}^n \otimes \mathbb{C}^n \otimes \mathbb{C}^n)^{\oplus m}$$

Thm [Dedersen, M.]

$$\beta(SL_{3n}^{\times 3}, (\mathbb{C}^{3n} \otimes \mathbb{C}^{3n} \otimes \mathbb{C}^{3n})^{\oplus 5}) \geq 4^n - 1$$

Thm [Dedersen, M.]

$$SL_{3n} \curvearrowright (\text{Sym}^3(\mathbb{C}^{3n}))^{\oplus 4}$$

\longrightarrow deg 3 polys in $3n$ variables.

$\beta \geq \frac{2}{3}(4^n - 1).$

 \longrightarrow indicate a pt for.

$$\mathbb{C}^{3n} \rightarrow \text{basis } \left\{ \begin{array}{l} x_1 \dots x_n \\ y_1 \dots y_n \\ z_1 \dots z_n \end{array} \right\}$$

$\text{Sym}^3 \mathbb{C}^{3n} = \text{deg } 3 \text{ polys in } \{x_i, y_i, z_i\}_{1 \leq i \leq n}$

\curvearrowright
 SL_{3n}

Idea: Lifting l.bds from torus actions.

GROSSMANS PRINCIPLE

$$G \curvearrowright V, W$$

$$v \in V$$

closed orbit
($G \cdot v$ is a closed subset)

$$G_v = \{g \in G \mid gv = v\} = H$$

$$\mathbb{C}[V \oplus W]^G \longrightarrow \mathbb{C}[W]^H$$

Cor: $\beta(G, V \oplus W) \geq \beta(H, W)$.

$$V \times W // G \longleftarrow G \cdot v \times W // G \begin{matrix} \parallel \\ W // H \end{matrix} \left[G \cdot v \approx G/H \right]$$

w_1, w_2 in the same H -orbit



(v, w_1) & (v, w_2) in the same G -orbit.

$G \curvearrowright V, W$

$v \in V$ closed orbit

$$G_v = \{g \in G \mid gv = v\}$$

$$\beta(G, v \oplus W) \geq \beta(H, W).$$

Steps:

- ① Identify a $v \in V$ with closed orbit
- ② Try and ensure H is a torus
- ③ Ability to prove 1-bd for H -action.

Challenge: How to prove that $v \in V$ has a closed orbit.

- ④ Getting all the steps to work at once.

Torus actions:

$$T = (\mathbb{C}^*)^k \curvearrowright V = \mathbb{C}^m \rightarrow \text{coordinates are } x_1, \dots, x_m$$

weight matrix

$$W = \begin{bmatrix} | & | \\ w_1 & w_m \\ | & | \end{bmatrix}$$

$$w_i \in \mathbb{Z}^k.$$

Invariant polynomials

$$\text{Ker}(W) \cap \mathbb{Z}_{\geq 0}^m$$

Inv. monomial

$$(a_1, \dots, a_m)$$

$$x_1^{a_1} \dots x_m^{a_m}$$

$$\sum a_i w_i = 0$$

Semi-group.

Generators \longleftrightarrow generators.

Eg: Rig things so that

$$W = \begin{bmatrix} 1 & -2 & & \\ & 1 & -2 & \\ & & 1 & -2 \\ & & & \ddots \end{bmatrix}$$

$$\text{Ker}(W) = \langle (8, 4, 2, 1) \rangle$$

Extending pattern

$$(2^n, 2^{n-1}, 2^{n-2}, \dots, 1)$$

$$\begin{matrix} & & & & \downarrow \text{corr. monomial} \\ x_1^{2^n} & x_2^{2^{n-1}} & \dots & x_n \end{matrix}$$

Hard part: Closed Orbit.

$$\beta(SL_{3n}, \text{Sym}^3(\mathbb{C}^{3n})^{\oplus 4}) \supseteq \frac{2}{3}(4^n - 1)$$

$$G = SL_{3n}, \quad V = \text{Sym}^3(\mathbb{C}^{3n})^{\oplus 3}$$

$$W = \text{Sym}^3(\mathbb{C}^{3n})$$

$$v = \left(\sum_i x_i^2 z_i, \sum_i y_i^2 z_i, \sum_i x_i y_i z_i \right)$$

$$\text{Sym}^3 \mathbb{C}^{3n}$$

① Orbit of v is closed.

② $T = (\mathbb{C}^*)^n \subset \mathbb{C}^{3n}$

$$(t_1, \dots, t_n) \cdot x_i = t_i x_i$$

$$y_i = t_i y_i$$

$$z_i = t_i^{-2} z_i$$

Stabilizer of v
i.e. $G_v = H$

$$\beta: (\mathbb{C}^*)^n \rightarrow SL_{3n}$$

$$\text{Im}(\beta) = G_v$$

$$\textcircled{3} \beta(T=H, \text{Sym}^3(\mathbb{C}^{3n})) \supseteq \frac{2}{3}(4^n - 1)$$

Why is the SL_{3n} -orbit of
 $v = (\sum x_i^2 z_i, \sum y_i^2 z_i, \sum \underbrace{2}_{} x_i y_i z_i)$ closed?

Idea: Moment map μ replace with numbers that have abs. value 1.

Kempf - Ness:

$\mu(v) = 0 \implies$ Orbit of v is closed.

\Downarrow
 v is a point of min. norm in the orbit $\} \rightarrow$ Such a point of min. norm exists if and only if the orbit is closed

Check that $\mu(v) = 0$.



How exactly to check $\mu(v) = 0$.

$$\mu: V \longrightarrow \text{Lie}(G)^*$$

$$v \longmapsto (X \longrightarrow \langle Xv, v \rangle)$$

Dadok-Kac Criterion:

$V = \bigoplus_{\lambda \in I} V_{\lambda}$ wt. space decomposition

$$v \in V \quad \text{supp}(v) = \{ \lambda \mid v_{\lambda} \neq 0 \}$$

"
 $(v_{\lambda})_{\lambda \in I}$

① $\text{conv}(\text{supp}(v)) \ni 0$ in the interior

② $\alpha, \beta \in \text{supp}(v) \Rightarrow \alpha - \beta$ is not a root.

① + ② $\Rightarrow \mu(v) = 0$.

Baty version $SL_3 \hookrightarrow \text{Sym}^3 \mathbb{C}^3$

$$v = x^3 + y^3 + z^3$$

$$\text{Supp}(v) = \left\{ (3, 0, 0), (0, 3, 0), (0, 0, 3) \right\} \left\{ \begin{array}{l} \text{modulo} \\ \langle (1, 1, 1) \rangle \end{array} \right\}$$

① $\text{Conv}(\text{supp}(v)) \ni 0$ in the interior
a day

② Roots are $\epsilon_i - \epsilon_j$

$$\left\{ \pm (1, -1, 0), \pm (1, 0, -1), \pm (0, 1, -1) \right\}$$

So ② is fine.

$$\begin{aligned} \underline{Q_n}: h_3(x, y, z) = & x^3 + y^3 + z^3 + x^2y + xy^2 \\ & + xz^2 + x^2z + y^2z + yz^2 \\ & + xyz \end{aligned}$$

Is the SL_3 -orbit closed?

SL_n actions (polynomial)

→ deg d poly actions

$d \geq 3 \rightarrow$ exp. bds

$d = 2 \rightarrow ?$

$d = 1 \rightarrow$ trivial

→ $SL_n \subset \text{Mat}_{n,n}^{n^2}$

$$g \cdot (A_1, \dots, A_{n^2}) = (gA_1g^{-1}, \dots, gA_{n^2}g^{-1})$$

Qn: Is the deg. bd for this action polynomial or exponential?